

REED-SOLOMON DECODER FOR PROCESSING (M) OR (2M) BIT DATA, AND DECODING METHOD THEREFOR

BACKGROUND OF THE INVENTION

Field of the Invention

The present invention relates to a Reed-Solomon (RS) decoder and a decoding method therefor, and more particularly to an RS decoder for processing m or (2m) bit data by using two RS cores in parallel, and a decoding method therefor. The present application is based on Korean Patent Application 2001-5970, filed February 7, 2001, which is incorporated herein by reference.

Description of the Related Art

In order to detect and correct an error during transmission or reproduction of data, a general digital communication and storing system employs an error correction code method in which transmission data is added with redundancy for error correction. The error correction code method is divided into a block code method and a tree code method. Here, block codes are divided into linear codes and nonlinear codes. The linear codes are divided into cyclic codes and non-cyclic codes. The cyclic code includes BCH code. The BCH code that forms code words in symbol units is the Reed-Solomon code (hereinafter, referred to as 'RS code'). For a code word of a certain correction capacity, the cyclic code or linear code requires more parity than the RS code. Accordingly, the RS codes of the cyclic code

method are relatively widely used in the digital communication and storing system when transmitting or storing a large amount of data. The RS code is a symbol-unit error correction code for detecting an error location and an error value from reception data according to a series of operation processes. In addition, the RS code has an excellent error correction ability for random and burst errors in a channel, and easily selects codes for obtaining a desired error rate. Especially, when used as a product code for generating an inner code and an outer code for the same data, the RS code for the digital communication and storing system transmitting or storing a large amount of data has an excellent correction ability.

A conventional RS decoder for processing (m) bit data will now be explained with reference to FIGS. 1 and 2.

The conventional RS decoder 100 includes an RS core 110, a control part 130 and a storing part 150.

The RS core 110 includes an eraser location polynomial calculation part (ERALCAL) 111, a syndrome polynomial calculation part (SYNDCAL) 113, an errata location polynomial calculation part (ERTLCAL) 115, and an error value calculation part (ERTVCAL) 117.

Here, eraser is defined as an error whose location can be detected.

Errata is the common name for an error or the eraser.

The control unit 130 connects the RS core 110 and the storing part 150, and generates a control signal CCON for controlling the RS core 110.

The storing unit 150 outputs to the control part 130 a data enable signal ACSEN for enabling the RS core 110 to access sequentially-inputted (m) bit data IDATA. According to the data enable signal ACSEN, the control unit 130 outputs a block offset address BADR and a block control signal BCON to the storing unit 150. The RS core 110 performs various operations by using the (m) bit data IDATA and an eraser flag ERAFLAG from the storing unit 150. The eraser flag ERAFLAG is an error flag calculated in the preceding step.

The SYNDICAL 113 calculates a syndrome polynomial from the sequentially-inputted (m) bit data IDATA.

The ERALCAL 111 calculates an eraser location polynomial having a root of an eraser location from the sequentially-inputted eraser flags ERAFLAG of the preceding step.

The ERTLICAL 115 calculates an errata location polynomial having a root of an errata location from the calculated syndrome polynomial and eraser location polynomial.

The ERTVCAL 117 calculates an error location ELOC and an error value EVAL from the calculated errata location polynomial and syndrome polynomial. The resultant error location ELOC, error value EVAL, error flag EFLAG1 and control signal STATUS are outputted to the control unit 130.

The control unit 130 decodes the data and corrects the error by using the output signals from the RS core 110, and outputs the error flag EFLAG2

and the decoded (m) bit data ODATA to the storing unit 150.

A conventional method for correcting an error of the RS decoder includes an inner code correction step and an outer code correction step.

The inner code correction step will be explained below. When a bit stream is inputted from a channel to an RS decoder 100, the bit stream is converted into (m) bit data symbols, and stored in the storing unit 150. This (m) bit data symbol is an error correction code block (ECC block) which satisfies the product of an inner code word length N1 and an outer code word length N2. When the data symbol is stored in the storing unit 150, the error correction starts. The RS core 110 sequentially reads the data from the storing unit 150 in (m) bit units, and the SYNDICAL 113 calculates the syndrome polynomial. After the syndrome polynomial is calculated, the error location polynomial is calculated. The error location and the error value are calculated from the calculated error location polynomial and syndrome polynomial. The error is corrected by reading an error symbol corresponding to the calculated error location, and then adding the error symbol to the calculated error value. Such a corrected error symbol is stored in the storing unit 150. If the error is detected precisely, '0' is stored as the eraser flag EFLAG2 value. Otherwise, '1' is stored as the eraser flag EFLAG2 value.

The outer code correction step is performed in the same order as the inner code correction step, except that the error flag location stored in the outer code corrections step is regarded as the eraser flag (ERAFLAG). Therefore,

the eraser flag ERAFLAG is read with the data in calculating the syndrome location polynomial, and the eraser location polynomial is calculated.

In FIG. 2, a lateral axis indicates time and a longitudinal axis indicates a procedure. In time period 0, a first procedure (proc1) is performed so that the eraser location polynomial and the syndrome polynomial are calculated by reading the eraser flag ERAFLAG and the first (m) bit data IDATA. In time period 1, the first procedure (proc1) is performed on the second (m) bit data, and a second procedure (proc2) for calculating the errata location polynomial is performed on the first (m) bit data. In time period 2, the first procedure (proc1) is performed on the third (m) bit data, the second procedure (proc2) for calculating the errata location polynomial is performed on the second (m) bit data, and a third procedure (proc3) for calculating the error location and the error value of the first (m) bit data, and updating the storing unit 150 is performed. That is, in the parallel processing method of a pipeline structure, while only the first procedure (proc1) is performed in time period 0, the first to third procedures (proc1-proc3) are simultaneously performed in time period 't'.

Even with the parallel processing of pipelined structure, however, the conventional decoder using the single RS core cannot satisfy the demand of a high-speed digital communication and storing system. In order to process the data suitable for the high-speed digital communication and storing system, the conventional RS decoder must have faster operation clocks. However, the

RS decoder, calculating a large amount of data, has limited operation clocks. In addition, when the speed of the operation clocks is increased to process the data at a high speed, an interface between the storing part and the RS core cannot be stabilized.

5 SUMMARY OF THE INVENTION

Accordingly, a primary object of the present invention is to provide an RS decoder for processing (m) or (2m) bit data which can satisfy the demand of a high-speed data communication and storing system.

Another object of the present invention is to provide an RS decoder for
10 processing (m) or (2m) bit data which can satisfy the demand of a high-speed data communication and storing system, without increasing the speed of the operation clocks.

Yet another object of the present invention is to provide an RS decoder for processing (m) or (2m) bit data which can be used for a complicated error
15 correction system.

An object of the present invention is to provide a method for an RS decoder to process (m) or (2m) bit data which satisfies the demand of a high-speed data communication and storing system.

Another object of the present invention is to provide a method for an
20 RS decoder to process (m) or (2m) bit data which can satisfy the demand of a high-speed data communication and storing system, without increasing the speed of the operation clocks.

Another object of the present invention is to provide a method for an RS decoder to process (m) or (2m) bit data which can be used for a complicated error correction system.

In order to achieve the above-described objects of the present invention,
5 there is provided an RS decoder including a storing part; a calculation part for calculating an error location and an error value from (2m) bit data from the storing part; and a control part for correcting an error of the data from the calculation part according to the error location and the error value, and controlling the calculation part to output a decoded signal.

10 In one embodiment of the present invention, the calculation part includes an eraser location polynomial calculation part for calculating an eraser location polynomial from an eraser flag from the storing part; a first syndrome polynomial calculation part for calculating a first syndrome polynomial from the data read from the storing part; a second syndrome
15 polynomial calculation part for calculating a second syndrome polynomial from the data read from the storing part; a first errata location polynomial calculation part for calculating a first errata location polynomial from the calculated eraser location polynomial and first syndrome polynomial, and outputting the delayed first syndrome polynomial; a first error location/value
20 calculation part for calculating a first error flag, a first error location and a first error value from the first errata location polynomial and the delayed first syndrome polynomial; a second errata location polynomial calculation part for

calculating a second errata location polynomial from the calculated eraser location polynomial and second syndrome polynomial, and outputting the delayed second syndrome polynomial; and a second error location/value calculation part for calculating a second error flag, a second error location and
5 a second error value from the second errata location polynomial and the delayed second syndrome polynomial.

In another embodiment of the present invention, the calculation part includes a first RS core for calculating a first error location and a first error value from the data read from the storing part; and a second RS core for
10 calculating a second error location and a second error value from the data read from the storing part.

In order to operate the RS decoder in (m) bit units, an (m) bit correction mode is set, in which the (m) bit data is stored in an up (m) bit memory of the storing part, and the second RS core is disabled not to access
15 the storing part.

The first RS core includes an eraser location polynomial calculation part for calculating an eraser location polynomial from an eraser flag read from the storing part; a first syndrome polynomial calculation part for calculating a first syndrome polynomial from the data read from the storing part; a first
20 errata location polynomial calculation part for calculating a first errata location polynomial from the calculated eraser location polynomial and first syndrome polynomial, and outputting the delayed first syndrome polynomial; and a first

error location/value calculation part for calculating a first error flag, a first error location and a first error value from the first errata location polynomial and the delayed first syndrome polynomial.

Preferably, the first syndrome polynomial calculation part satisfies $S_j = \alpha^j(S_{j-1}\alpha^j + UM) + DM$ when $(2m)$ bit data is inputted, and satisfies $S_j = S_{j-1}\alpha^j + UM$ when (m) bit data is inputted. Where S_j indicates a current state syndrome polynomial, S_{j-1} is a preceding state syndrome polynomial, α^j is a root of a generated polynomial, UM is up (m) bits of the $(2m)$ bit data, and DM is down (m) bits of the $(2m)$ bit data.

The second RS core includes a second syndrome polynomial calculation part for calculating a second syndrome polynomial from the data read from the storing part; a second errata location polynomial calculation part for calculating a second errata location polynomial from the calculated eraser location polynomial and second syndrome polynomial, and outputting the delayed second syndrome polynomial; and a second error location/value calculation part for calculating a second error flag, a second error location and a second error value from the second errata location polynomial and the delayed second syndrome polynomial.

The second syndrome polynomial calculation part satisfies $S_j = \alpha^j(S_{j-1}\alpha^j + UM) + DM$ when $(2m)$ bit data is inputted, and satisfies $S_j = S_{j-1}\alpha^j + DM$ when (m) bit data is inputted. Where S_j indicates a current state syndrome polynomial, S_{j-1} is a preceding state syndrome polynomial, α^j is a

root of a generated polynomial, UM is up (m) bits of the (2m) bit data, and DM is down (m) bits of the (2m) bit data.

In another embodiment of the present invention, the RS decoder
5 includes a storing part for storing (2m) bit data; a main control part for
controlling the storing part and overall operation of the decoder; a first RS core
for calculating a first error location and a first error value from the data read
from the storing part; a first RS core control part for controlling the first RS
core under the control of the main control part; a second RS core for
10 calculating a second error location and a second error value from the data read
from the storing part; and a second RS core control part for controlling the
second RS core under the control of the main control part.

In one embodiment of the present invention, an RS decoding method
includes the steps of reading the data to be decoded and an eraser flag read
15 from the storing part; calculating an error location and an error value from the
read data; and correcting an error of the data according to the calculated error
location and error value, and decoding the data.

Preferably, the data is read in (2m) bit units in the data reading step.

The data reading step includes an eraser location polynomial
20 calculation step for calculating an eraser location polynomial from the eraser
flag from the storing part; a first syndrome polynomial calculation step for
calculating a first syndrome polynomial from the read data; a second syndrome

polynomial calculation step for calculating a second syndrome polynomial from the read data; a first errata location polynomial calculation step for calculating a first errata location polynomial from the calculated eraser location polynomial and first syndrome polynomial, and outputting the delayed first syndrome polynomial; a first error location/value calculation step for calculating a first error flag, a first error location and a first error value from the first errata location polynomial and the delayed first syndrome polynomial; a second errata location polynomial calculation step for calculating a second errata location polynomial from the calculated eraser location polynomial and second syndrome polynomial, and outputting the delayed second syndrome polynomial; and a second error location/value calculation step for calculating a second error flag, a second error location and a second error value from the second errata location polynomial and the delayed second syndrome polynomial.

15 The calculation step includes a first calculation step for calculating a first error location and a first error value from the read data; and a second calculation step for calculating a second error location and a second error value from the read data.

20 The first calculation step includes an eraser location polynomial calculation step for calculating an eraser location polynomial from the eraser flag; a first syndrome polynomial calculation step for calculating a first syndrome polynomial from the read data; a first errata location polynomial

calculation step for calculating a first errata location polynomial from the
calculated eraser location polynomial and first syndrome polynomial, and
outputting the delayed first syndrome polynomial; and a first error
location/value calculation step for calculating a first error flag, a first error
5 location and a first error value from the first errata location polynomial and the
delayed first syndrome polynomial.

Preferably, the first syndrome polynomial calculation step satisfies $S_j = \alpha^j(S_{j-1}\alpha^j + UM) + DM$ when $(2m)$ bit data is inputted, and satisfies $S_j = S_{j-1}\alpha^j + UM$ when (m) bit data is inputted. S_j indicates a current state
10 syndrome polynomial, S_{j-1} is a preceding state syndrome polynomial, α^j is a
root of a generated polynomial, UM is up (m) bits of the $(2m)$ bit data, and
 DM is down (m) bits of the $(2m)$ bit data.

The second calculation step includes a second syndrome polynomial
calculation step for calculating a second syndrome polynomial from the read
15 data; a second errata location polynomial calculation step for calculating a
second errata location polynomial from the calculated eraser location
polynomial and second syndrome polynomial, and outputting the delayed
second syndrome polynomial; and a second error location/value calculation
step for calculating a second error flag, a second error location and a second
20 error value from the second errata location polynomial and the delayed second
syndrome polynomial.

The second syndrome polynomial calculation step satisfies $S_j =$

$\alpha^j(S_{j-1}\alpha^j+UM)+DM$ when (2m) bit data is inputted, and satisfies $S_j = S_{j-1}\alpha^j+DM$ when (m) bit data is inputted. S_j indicates a current state syndrome polynomial, S_{j-1} is a preceding state syndrome polynomial, α^j is a root of a generated polynomial, UM is up (m) bits of the (2m) bit data, and
5 DM is down (m) bits of the (2m) bit data.

According to the present invention, an inner code correction method of an RS production code includes the steps of calculating a first syndrome polynomial from inner code words received in (2m) bit units; calculating a second syndrome polynomial from inner code words received in (2m) bit
10 units; calculating first and second errata location polynomials from the calculated first and second syndrome polynomials and an eraser location polynomial; and calculating first and second error values and first and second error locations according to the first and second errata location polynomials and the first and second syndrome polynomials, wherein the errors are
15 corrected in (m) bit units. The order of the first and second syndrome polynomial calculation can be changed.

An outer code correction method of an RS production code according to the present invention includes the steps of calculating a first syndrome polynomial from up (m) bits of read (2m) bit outer codes and a second
20 syndrome polynomial from down (m) bits of read (2m) bit outer codes, and simultaneously calculating an eraser location polynomial by reading an eraser flag; calculating first and second errata location polynomials from the first and

second syndrome polynomials and the eraser location polynomial; and calculating an error value and an error location from the first and second errata location polynomials and the first and second syndrome polynomials, the errors being alternately corrected in (m) bit units.

- 5 According to the present invention, the RS decoder for processing (m) or (2m) bit data and the decoding method therefor can process the data at a high speed, by using the RS cores in parallel. Therefore, it is unnecessary to increase a speed of operation clocks, thus stabilizing an interface between the data storing part and the RS core, and increasing reliability of the system.
- 10 Moreover, the syndrome polynomial calculation part of the RS core selectively processes (m) or (2m) bit data. As a result, the present invention can be employed for CD/DVD data processing of a complicated error correction system.

BRIEF DESCRIPTION OF THE DRAWINGS

- 15 A more complete appreciation of the invention, and many of the attendant advantages thereof, will be readily apparent as the same becomes better understood by reference to the following detailed description when considered in conjunction with the accompanying drawings in which like reference symbols indicate the same or similar components, wherein:

- 20 FIG. 1 is a block diagram illustrating a conventional RS decoder for processing (m) bit data;

FIG. 2 is a timing diagram showing data process of the RS decoder in

FIG. 1 in each time period;

FIG. 3 illustrates data structure of an RS production code;

FIG. 4 is a block diagram illustrating an RS decoder for processing (m) or (2m) bit data in accordance with the present invention;

5 FIG. 5 is a timing diagram showing inner code data processing of the RS decoder in FIG. 4 in each time period;

FIG. 6 is a timing diagram showing outer code data processing of the RS decoder in FIG. 4 in each time period; and

10 FIG. 7 is a detailed structure diagram illustrating a syndrome calculation part of the RS decoder in FIG. 4.

Fig. 8 is a diagram illustrating a Reed-Soloman decoding method in accordance with the present invention.

Fig. 9 is a diagram illustrating a data reading method in accordance with the present invention.

15 Fig. 10 is a diagram illustrating a calculation method in accordance with the present invention.

Fig. 11 is a diagram illustrating a first and a second calculation method in accordance with the present invention.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENT

20 A Reed-Solomon (RS) decoder for processing (m) or (2m) bit data, and a decoding method therefor in accordance with the present invention will now be described with reference to the accompanying drawings.

For a reference, the character 'm' of the (m) or (2m) bit data is a unit that represents amount of data processing, for example, a number of bits constituting a symbol, which may be appropriately determined according to a data representation method. Hereinafter, the basic unit that represents amount of data processing will be represented by 'm'.

RS(N, K, d) is an RS code whose code word length is 'N', information word length is 'K', and minimum hamming distance is 'd'. One of the characteristics of the RS code is that the minimum hamming distance (d) is equal to (N-K+1), where 'N-K' is a parity number. When 'N-K' is 'R', 'R' is equal to (d-1). The formula $N-K=R=d-1$ is for a definition of the number of parity. According to the formula, for example, if an information word is 8 bits (K=8), the length of the signal word is 12 (N=12) and the number of the parity is 4 (R=N-K). When the number of symbols correctable by the RS code is 't', 't' is equal to $[(d-1)/2]$. In the code word including 'e' erasers, the RS code can correct 't' errors and 'e' erasers which satisfy $d \geq 2t+e+1$. The RS code for a digital communication and storing system for transmitting or storing a large amount of data generally uses a product code that encodes an inner code and an outer code generated with respect to the same data. FIG. 3 illustrates data structure of the RS production code RSPC. The RS production code includes data, an inner parity and an outer parity. Here, an inner code length is 'N1' and an outer code length is 'N2'.

FIG. 4 is a block diagram illustrating an RS decoder 200 for processing

(m) or (2m) bit data in accordance with a first embodiment of the present invention.

Referring to FIG. 4, the RS decoder 200 includes a storing unit 210, a main control unit 220 for controlling the storing unit 210 and overall operation of the decoder, a first RS core 230, a first RS core control unit 240, a second RS core 250 and a second RS core control unit 260.

The storing unit 210 can be a buffer or a memory for storing (2m) bit data.

The first RS core 230 calculates a first error location ELOC1 and a first error value EVAL1 from the data IDATA read from the storing unit 210.

The first RS core control unit 240 controls the first RS core 230 under the control of the main control unit 220.

The second RS core 250 calculates a second error location ELOC2 and a second error value EVAL2 from the data IDATA read from the storing unit 210.

The second RS core control unit 260 controls the second RS core 250 under the control of the main control unit 220.

The first RS core 230 includes an eraser location polynomial calculation part ERALCAL 231, a first syndrome polynomial calculation part SYNDCAL1 233, a first errata location polynomial calculation part ERTLAL1 235 and a first error location/value calculation part ERTVCAL1 237.

The ERALCAL 231 calculates an eraser location polynomial ① from an eraser flag ERAFLAG read from the storing unit 210. The eraser flag ERAFLAG is an error flag (EFLAG) calculated in a preceding step. The eraser is data that has information about a position of an error but not about a value of the error. Referring to Fig. 3, the horizontal axis indicates an internal signal correction direction. If there are data which are not corrected even after the process of correcting the internal signal, it means that an error has occurred in the data along the horizontal axis, and accordingly, one error flag is stored. And, this stored error flag is, from the horizontal axis point of view, seen as an eraser flag because it indicates a position where the error occurs.

The SYNDCAL1 233 calculates a first syndrome polynomial ② from the data IDATA read from the storing unit 210.

The ERTLAL1 235 calculates a first errata location polynomial ③ from the calculated eraser location polynomial ① and first syndrome polynomial ②, and outputs the delayed first syndrome polynomial ④.

The ERTVCAL1 237 calculates a first error flag EFLAG1, a first error location ELOC1 and a first error value EVAL1 from the first errata location polynomial ③ and the delayed first syndrome polynomial ④.

The second RS core 250 includes a second syndrome polynomial calculation part SYNDCAL2 253, a second errata location polynomial

calculation part ERTLAL2 255 and a second error location/value calculation part ERTVAL2 257.

The SYNDAL2 253 calculates a second syndrome polynomial ⑤ from the data IDATA read from the storing unit 210.

5 The ERTLAL2 255 calculates a second errata location polynomial ⑥ from the calculated eraser location polynomial ① and second syndrome polynomial ⑤, and outputs the delayed second syndrome polynomial ⑦.

10 The ERTVAL2 257 calculates a second error flag EFLAG2, a second error location ELOC2 and a second error value EVAL2 from the second errata location polynomial ⑥ and the delayed second syndrome polynomial ⑦.

The outputs ⑧ and ⑨ are values including the error values that are calculated from the first RS core 230 and the second RS core 250, respectively, and that respectively flow from the first RS core control unit 220 and from the second RS core control unit 260 to the main control unit 220.

15 The operation of the (2m) bit input/output decoder for the inner code and the outer code will now be described with reference to FIGS. 3 through 6. In FIGS. 5 and 6, a single straight line indicates a predetermined operation in progress, and double straight lines indicate an update operation.

20 The inner code correction operation will now be explained with reference to FIGS. 4 and 5. In time period 0, the first RS core 230 receives data in (2m) bit units, and the SYNDAL1 233 calculates the first syndrome polynomial in (2m) bit units (P11). Thereafter, the second RS core 250

receives data in $(2m)$ bit units, and the SYNDCAL2 253 calculates the second syndrome polynomial in $(2m)$ bit units (P21). Here, the procedures (P11) and (P12) may be carried out in an inverse order. The syndrome calculation is performed in $(2m)$ bit units. Compared with the conventional (m) bit unit
5 syndrome calculation part, the calculation time is reduced into a half. That is, the calculation speed of the syndrome polynomial is increased by two times.

In time period 1, the ERTLAL1 235 and the ERTLAL2 255 receive the calculated first and second syndrome polynomials ② and ⑤ and eraser location polynomial ①, and calculate the first and second errata location
10 polynomials ③ and ⑥(P12, P22). While the errata location polynomial is performed on the first $(2m)$ bit data, the first and second RS cores 230, 250 receive the second $(2m)$ bit data, and calculate the first and second syndrome polynomials ② and ⑤(P11, P21).

In time period 2, the error value EVAL1 and the error location ELOC1
15 of the first $(2m)$ bit data is calculated, and the error correction operation is performed. Here, the error correction operation is performed in (m) bit units.

Accordingly, the first RS core 230 and the second RS core 250 are alternately enabled under the control of the main control unit 220, to correct the error and update the data (P13, P23). At the same time, the first and second errata
20 location polynomials ③ and ⑥ are calculated on the second $(2m)$ bit data (P12, P22), and the first and second syndrome polynomials ② and ⑤ are

calculated on the third (2m) bit data (P11, P21). As described above, the syndrome polynomial calculation ② and ⑤ (P11, P21), the errata location calculation ③ and ⑥ (P12, P22) and the error correction and data updating operations (P13, P23) are simultaneously performed in a predetermined time period after time period 2. When the error correction operation is finished with respect to the inner code, the identical procedure is performed on the two succeeding inner code words. Therefore, the inner code correction operation is performed N1/2 times.

The outer code correction operation will now be described with reference to FIGS. 4 and 6.

In time period 0, a (2m) bit unit, namely, the two outer codes, is read. The first RS core 230 receives up (m) (UM) bits and the second RS core 250 receives down (m) (DM) bits, thereby calculating the first and second syndrome polynomials ② and ⑤ at the same time. In addition, the ERALCAL 231 receives the eraser flag ERAFLAG, and calculates the eraser location polynomial ① (P15, P25). Accordingly, the syndrome polynomial ② and ⑤ and the eraser location polynomial ① for the two outer codes are obtained during the general (m) bit unit syndrome calculation time.

In time period 1, the ERTLAL1 235 and the ERTLAL2 255 receive the calculated first and second syndrome polynomials ② and ⑤ and eraser location polynomial ①, and calculate the first and second errata location

polynomials ③ and ⑥ (P17, P27). While the errata location polynomial calculation is performed on the first (2m) bit outer code, the first and second RS cores 230, 250 receive the second (2m) bit outer code, and calculate the first and second syndrome polynomials ② and ⑤ (P15, P25).

5 In time period 2, the error value EVAL1 and the error location ELOC1 of the first (2m) bit outer code are calculated, and the error correction operation is performed. Here, the error correction operation must be performed in (m) bit units. Accordingly, the first RS core 230 and the second RS core 250 are alternately enabled under the control of the main control unit
10 220, to correct the error and update the data (P19, P29). At the same time, the first and second errata location polynomials ③ and ⑥ are calculated on the second (2m) bit outer code (P17, P27), and the first and second syndrome polynomials ② and ⑤ are calculated on the third (2m) bit outer codes (P15, P25). As described above, the syndrome polynomial calculation ② and ⑤
15 (P15, P25), the errata location calculation ③ and ⑥ (P17, P27) and the error correction and data updating operations (P19, P29) are simultaneously performed in a predetermined time period after time period 2. When the error correction operation is finished, the identical procedure is performed on the two succeeding outer code words. Therefore, the outer code correction
20 operation is performed $N/2$ times.

The operation of the (2m) bit unit RS decoder was explained above.

In order to operate the (m) bit unit RS decoder, an (m) bit correction mode is set. In the (m) bit correction mode, the (m) bit data is stored in an up (m) (UM) bit memory of the storing unit 210, and the second RS core 250 is disabled so that the second RS core 250 cannot access the storing unit 210. Accordingly, the (m) bit unit RS decoder operates in the same way as the conventional (m) bit decoder does.

As illustrated in FIG. 7, the SYNDICAL 233 includes a first syndrome storing part 233a, a first multiplier 233b, a first adder 233c, a first (m) bit multiplexer 233d, a second multiplier 233e, a second (m) bit multiplexer 233f and a second adder 233g.

The first syndrome storing part 233a temporarily stores a calculation result of the first syndrome polynomial.

The first multiplier 233b multiplies the syndrome polynomial from the first syndrome storing part 233a by a root α^j ($j=0, 1, \dots, N-K-1$) of the generated polynomial.

The first adder 233c adds the output from the first multiplier 233b to up (m) (UM) bits of the input data.

The first (m) bit multiplexer 233d respectively outputs 1 or α^j according to the (m) or (2m) bit mode.

The second multiplier 233e multiplies the output from the first adder 233c by the output from the first (m) bit multiplexer 233d.

The second (m) bit multiplexer 233f respectively outputs 0 or down

(m) (DM) bits of the input data according to (m) or (2m) bit mode.

The second adder 233g adds the output from the second multiplier 233e to the output from the second (m) bit multiplexer 233f. The first syndrome storing part 233a temporarily stores and outputs the output from the
5 second adder 233g.

The SYNDCAL2 253 includes a second syndrome storing part 253a, a third (m) bit multiplexer 253b, a third multiplier 253c, a fourth (m) bit multiplexer 253d, a third adder 253e, a fourth multiplier 253f and a fourth adder 253g.

10 The second syndrome storing part 253a temporarily stores a calculation result of the second syndrome polynomial.

The third (m) bit multiplexer 253b respectively outputs 1 or α^j according to the (m) or (2m) bit mode.

The third multiplier 253c multiplies the output from the second
15 syndrome storing part 253a by the output from the third (m) bit multiplexer 253b.

The fourth (m) bit multiplexer 253d respectively outputs 0 or up (m) (UM) bits of the input data according to (m) or (2m) bit mode.

The third adder 253e adds the output from the third multiplier 253c to
20 the output from the fourth (m) bit multiplexer 253d.

The fourth multiplier 253f multiplies the output from the third adder 253e by a root α^j ($j=0, 1, \dots, N-K-1$) of the generated polynomial.

The fourth adder 253g adds the output from the fourth multiplier 253f to the down (m) (DM) bits of the input data. The second syndrome storing part 253a stores and outputs the output from the fourth adder 253g.

As illustrated in FIG. 8, the high speed RS decoding method in accordance with the present invention includes the steps of reading the data to be decoded and an eraser flag (step S1); calculating an error location and an error value from the read data (step S2); and correcting an error of the data according to the calculated error location and error value, and decoding the data (step S3).

Preferably, the data is read in (2m) bit units in the data reading step (S1).

As illustrated in FIG. 9, The data reading step (S1) includes a read eraser flag step (S1.1), an eraser location polynomial calculation step (step S11), a first syndrome polynomial calculation step (step S12), a second syndrome polynomial calculation step (step S13), a first errata location polynomial calculation step (step S14), a delayed first syndrome polynomial output step (S14.1), a first error location/value calculation step (step S15), a first error flag calculation output step (S15.1), a first error location calculation output step (S15.2), a first error value calculation output step (S15.3), a second errata location polynomial calculation step (step S16), a delayed second syndrome polynomial output step (S16.1), a second error location/value calculation step (step S17), a second error flag calculation output step (S17.1),

a second error location calculation output step (S17.2), and a second error value calculation output step (S17.3).

In the eraser location polynomial calculation step (S11), an eraser location polynomial is calculated by using the read eraser flag (S1.1).

5 In the first syndrome polynomial calculation step (S12), a first syndrome polynomial is calculated from the read data.

In the second syndrome polynomial calculation step (S13), a second syndrome polynomial is calculated from the read data.

10 In the first errata location polynomial calculation step (S14), a first errata location polynomial is calculated from the calculated eraser location polynomial and first syndrome polynomial, and the delayed first syndrome polynomial is outputted (S14.1).

15 In the first error location/value calculation step (S15), a first error flag, a first error location and a first error value are calculated from the first errata location polynomial and the delayed first syndrome polynomial (S14.1). Then, the first error flag is output (S15.1), the first error location is output (S15.2), and the first error value is output (S15.3).

20 In the second errata location polynomial calculation step (S16), a second errata location polynomial is calculated from the calculated eraser location polynomial and second syndrome polynomial, and the delayed second syndrome polynomial is outputted (S16.1).

In the second error location/value calculation step (S17), a second error

flag, a second error location and a second error value are calculated from the second errata location polynomial and the delayed second syndrome polynomial (S16.1). Then, the second error flag is output (S17.1), the second error location is output (S17.2), and the second error value is output (S17.3).

- 5 The calculation step (S2) is illustrated in FIG. 10. It includes a first calculation step (step S21) and a second calculation step (step S22).

As illustrated in FIG. 11, the first calculation step (S21) includes a first error location and a first error value calculated from the read data.

- 10 In the second calculation step (S22), a second error location and a second error value are calculated from the read data.

- 15 The first calculation step (S21) includes an eraser location polynomial calculation step (step S211), a first syndrome polynomial calculation step (step S212), a first errata location polynomial calculation step (step S213), a delayed first syndrome polynomial output step (S213.1), a first error location/value calculation step (step S214), a first error flag calculation output step (S214.1), a first error location calculation output step (S214.2), and a first error value calculation output step (S214.3).

In the eraser location polynomial calculation step (S211), an eraser location polynomial is calculated from the read eraser flag.

- 20 In the first syndrome polynomial calculation step (S212), a first syndrome polynomial is calculated from the read data. The first syndrome polynomial calculation step (S212) satisfies $S_j = \alpha^j(S_{j-1}\alpha^j + UM) + DM$ when

(2m) bit data is inputted, and satisfies $S_j = S_{j-1}\alpha^j + UM$ when (m) bit data is inputted. Here, S_j indicates a current state syndrome polynomial, S_{j-1} is a preceding state syndrome polynomial, α^j is a root of the generated polynomial, UM is up (m) bits of the (2m) bit data, and DM is down (m) bits of the (2m) bit data.

In the first errata location polynomial calculation step (S213), a first errata location polynomial is calculated from the calculated eraser location polynomial and first syndrome polynomial, and the delayed first syndrome polynomial is outputted (213.1).

In the first error location/value calculation step (S214), a first error flag, a first error location and a first error value are calculated from the first errata location polynomial and the delayed first syndrome polynomial (213.1). Then, the first error flag is output (S214.1), the first error location is output (S214.2), and the first error value is output (S214.3).

The second calculation step (S22) includes a second syndrome polynomial calculation step (step S221), a second errata location polynomial calculation step (step S222), a delayed second syndrome polynomial output step (S222.1), a second error location/value calculation step (step S223), a second error flag calculation output step (S223.1), a second error location calculation output step (S223.2), and a second error value calculation output step (S223.3).

In the second syndrome polynomial calculation step (S221), a second

syndrome polynomial is calculated from the read data. Preferably, the second syndrome polynomial calculation step (S221) satisfies $S_j = \alpha^j(S_{j-1}\alpha^j + UM) + DM$ when (2m) bit data is inputted, and satisfies $S_j = S_{j-1}\alpha^j + DM$ when (m) bit data is inputted.

5 In the second errata location polynomial calculation step (S222), a second errata location polynomial is calculated from the calculated eraser location polynomial and second syndrome polynomial, and the delayed second syndrome polynomial is outputted (222.1).

10 In the second error location/value calculation step (S223), a second error flag, a second error location and a second error value are calculated from the second errata location polynomial and the delayed second syndrome polynomial (222.1). Then, the second error flag is output (S223.1), the second error location is output (S223.2), and the second error value is output (S223.3).

15 According to the present invention, the RS decoder for processing (m) or (2m) bit data and the decoding method therefor can process the data at a high speed, by using the RS cores in parallel. Therefore, it is unnecessary to increase the speed of the operation clocks, thus stabilizing an interface between the storing part and the RS core, and increasing the reliability of the
20 system.

Moreover, the syndrome polynomial calculation part of the RS core selectively processes (m) or (2m) bit data. As a result, the present invention

can be employed for CD/DVD data processing of a complicated error correction system. For example, the present invention can be employed for a CD error correction system, which is a Cross Interleaved Reed-Solomon Code (CIRC) that needs to decode data in (m) bit units in order to perform a complicated interleave operation.

While the present invention has been particularly shown and described with reference to the preferred embodiments thereof, it will be understood by those skilled in the art that various changes in form and details may be effected therein without departing from the spirit and scope of the invention as defined by the appended claims.